

## A comparison between two collocations methods for linear polylocal problems - a Computer Algebra based approach

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### Abstract

*Consider the problem:*

$$\begin{aligned} -y''(t) + q(t)y(t) &= r(t), & t \in [a, b] \\ y(c) &= \alpha \\ y(d) &= \beta, & c, d \in (a, b). \end{aligned}$$

*The aim of this paper is to present two approximate solutions of this problem based on B-splines and first kind Chebyshev polynomials, respectively. The first solution uses a mesh based on Legendre points, while the second uses a Chebyshev-Lobatto mesh. Using computer algebra techniques and a Maple implementation, we obtain analytical expression of the approximations and give examples. Chebyshev method has a smaller error, but for large number of mesh points the B-spline method is faster and requires less memory.*

## 1 Introduction

Consider the problem:

$$-y''(t) + q(t)y(t) = r(t), \quad t \in [a, b] \tag{1}$$

$$y(d) = \alpha \tag{2}$$

$$y(e) = \beta, \quad d, e \in (a, b), d < e. \tag{3}$$

where  $q, r \in C[a, b]$ ,  $\alpha, \beta \in \mathbb{R}$ . This is not a two-point boundary value problem, since  $d, e \in (a, b)$ .

If the solution of the two-point boundary value problem

$$\begin{aligned} -y''(t) + q(t)y(t) &= r(t), & t \in [d, e] \\ y(d) &= \alpha \\ y(e) &= \beta, \end{aligned} \tag{4}$$

exists and it is unique, then the requirement  $y \in C^2[a, b]$  assures the existence and the uniqueness of (1)+(2)+(3).

We have two initial value problems on  $[a, d]$  and  $[e, b]$ , respectively, and the existence and the uniqueness for (4) assure existence and uniqueness of these problems. It is possible to solve this problem by dividing it into the three above-mentioned problems and to solve each of these problem separately, but we are interested to a unitary approach that solve it as a whole.

In 1966, two researchers from Tiberiu Popoviciu Institute of Romanian Academy, Cluj-Napoca, *Dumitru Ripianu* and *Oleg Arama* published a paper on a polylocal problem, see [9].

## 2 Principles of the method

The implementation is inspired from [4, 5].

### 2.1 B-spline method

Our first method is based on collocation with nonuniform cubic B-splines [2, 10]. For properties of B-spline and basic algorithms see [5].

Consider the mesh (see [1])

$$\Delta : a = x_0 < x_1 < \dots < x_m < x_{m+1} = b, \quad (5)$$

and the step sizes

$$h_i := x_{i+1} - x_i, \quad i = 0, \dots, m.$$

Within each subinterval we insert  $k$  points

$$0 \leq \rho_1 < \rho_2 < \dots < \rho_k \leq 1,$$

which are the roots of the  $k$ th Legendre's orthogonal polynomial on  $[0, 1]$  [6, 8].

Finally, the mesh has the form

$$\xi_{i,j} := x_i + h_i \rho_j, \quad j = 1, \dots, k, \quad i = 0, \dots, m.$$

The number of mesh points is now  $N = (m + 1)k$ .

We shall choose the basis such that the following conditions hold:

- the solution verifies the differential equation (1) at  $\xi_{i,j}$ ;
- the solution verifies the conditions (2), (3).

We need a basis having  $N + 2$  cubic B-spline functions.

One rennumbers the points such that the first point is  $x_0$  and the last is  $x_{n+1}$ .

In order to impose the fulfillment of (1) at  $a$  and  $b$  we complete the mesh with points  $x_{-k}, x_{-k+1}, \dots, x_{-1}$  and  $x_{n+2}, x_{n+3}, \dots, x_{n+k+1}$ .

The form of solution is

$$y(t) = \sum_{i=-1}^{n+2} b_i B_i(t), \quad (6)$$

where  $B_i(t)$  is the B-spline with knots  $x_{k-2}, x_{k-1}, x_k, x_{k+1}, x_{k+2}$ .

The conditions on solution yield a linear system with  $n + 4$  equations and  $n + 4$  unknowns (the coefficients  $b_i$ ,  $i = -1, \dots, n + 2$ ).

The system matrix is banded with at most 4 nonzero elements on each line (3 nonzero at each mesh point and four at  $d$  and  $e$ ).

### 2.2 Chebyshev method

Our second method is based on first kind Chebyshev polynomials [6, 8]. We consider the mesh

$$\frac{b-a}{2} \cos \frac{k\pi}{n} + \frac{a+b}{2}, \quad k = 0, \dots, n \quad (7)$$

(the extremes of Chebyshev #1 polynomials, or equivalently the roots of Chebyshev #2 polynomials) completed with inner points  $c$  and  $d$ . The form of the solution is

$$y(t) = \sum_{i=0}^{n+1} c_i T_i(t); \quad (8)$$

where  $T_i(t)$  is the  $k$ -th degree first kind Chebyshev polynomial on interval  $[a, b]$ . As in the previous section, the fulfillment of (1), (2) and (3) leads us to a system of  $n + 2$  equations and  $n + 2$  unknowns (the coefficients  $c_i$ ,  $i = 0, \dots, n + 1$ ). This time the matrix is dense.

### 3 Maple implementation

We implement our ideas in Maple 10. For necessary details on Maple see [7]. Both methods return the approximation in analytic form.

#### 3.1 B-spline method

The basic functions are computed using the function `BSpline` of the package `CurveFitting`. The B-spline basis is obtained through Maple sequence

```
> S:=(x,u,k)->eval(BSpline(4,t,
> knots=[seq(u[i],i=k-2..k+2)]),
> t=x):
```

`S(x,u,k)` computes the cubic B-spline in variable  $x$ , with knots  $u[k-2], \dots, u[k+2]$ .

The procedure `genspline` computes the B-spline solution. It accepts the mesh  $x$ , the number of points  $n$ , the functions  $q$  and  $r$ , the points  $d$ ,  $e$  and the values at  $d$  and  $e$ ,  $\alpha$  and  $\beta$ , respectively. It returns the solution  $y$ , given by (6). The matrix of the system and the right-hand side vector are constructed element by element and the solution is computed using the function `LinearSolve` from `LinearAlgebra` package. This is a fast and flexible solution, and allows the selection of the solution method and gaining additional information, like condition number. Here is the Maple code.

```
> genspline:=proc(x,n,q,r,d,e,
> alpha,beta)
> local k, i, A, y, poze, pozd, ii,p,xe,xd, Y;
> global S, b;
> A:=Matrix(n+4,n+4); y:=Vector(n+4):
> b:=Vector(n+4):
> ii:=1;
> for i from 0 to n+1 do
> for k from max(i-1,-1) to i+1 do
> A[ii,k+2]:=(-eval(diff(S(t,x,k),
> t$2), t=x[i])+q(x[i])*
> eval(S(t,x,k),t=x[i]));
> end do:
> y[ii]:=r(x[i]);
> if (x[i]<d and x[i+1]>d) then
> ii:=ii+1; pozd:=ii; xd:=i;
> end if:
> if (x[i]<e and x[i+1]>e) then
> ii:=ii+1; poze:=ii; xe:=i;
> end if:
> ii:=ii+1;
> end do:
> p:=xd;
> for k from p-1 to p+2 do
> A[pozd,k+2]:=eval(S(t,x,k),t=d);
> end do;
> y[pozd]:=alpha;
> p:=xe;
> for k from p-1 to p+2 do
> A[poze,k+2]:=eval(S(t,x,k),t=e);
> end do;
> y[poze]:=beta;
> b:=LinearSolve(A,y);
> Y:=0:
> for k from -1 to n+2 do
> Y:=Y+b[k+2]*S(t,x,k):
> end do:
> return Y:
> end proc:
```

The procedure `genspline` accepts the mesh given in array form. The procedure `gendivLeg` generates the mesh as shown in Section 2.1. It calls the procedure `genpoints`. It computes the Legendre polynomial, solve it, and generates mesh points using an affine transform. The Legendre polynomials are generated using the `orthopoly` package, and their roots are obtained via `solve` function. Here is the code for `genpoints`:

```
> genpoints:=proc(a,b,N,k)
> local L,i,j,xu,xc,pol,pol2,sol,
> h,nL,x;
> L:=[a];
> h:=(b-a)/(N+1);
> xc:=a-h; pol:=P(k,t);
> pol2:=expand(subs(t=2*x-1,pol));
> sol:=fsolve(pol2);
> for i from 0 to N+2 do
> xu:=xc+h;
> for j from 1 to k do
> L:=[op(L),xc+(xu-xc)*sol[j]];
> end do;
> xc:=xu;
> end do;
> L:=[op(L),b];
> L:=sort(L);
> return L;
> end proc;
```

The code for `gendivLeg` closes the section.

```
> gendivLeg:=proc(a,b,n,k)
> local h,x,Y,L,nn,j;
> L:=genpoints(a,b,n,k);
> L:=convert(L,rational,exact);
> nn:=nops(L)-2*k;
> x:=Array(-k..nn+k-1,L);
> return x;
> end proc;
```

### 3.2 Chebyshev method

The Chebyshev polynomials are generated via the `orthopoly` package. The Maple sequence

```
> S:=(x,k,a,b)->T(k,
> ((b-a)*x+a+b)/2):
```

computes the  $k$ -th degree Chebyshev polynomial on interval  $[a, b]$ . The following Maple procedure `genceb` is the analogous of `genspline`. It uses `solve` to compute the Chebyshev coefficients.

```

> genceb:=proc(x,n,q,r,c0,d0,
> alpha,beta)
> local k, ecY, ecd, C, h, Y, c, a, b;
> global S;
> a:=x[0]; b:=x[n-1];
> Y:=0;
> for k from 0 to n+1 do
> Y:=Y+c[k]*S(t,k,a,b);
> end do;
> Y:=simplify(Y);
> ecY:=-diff(Y,t$2)+q(t)*Y=r(t):
> ecd:=Array(0..n+1);
> for k from 0 to n-1 do
> ecd[k]:=eval(ecY,t=x[k]):
> end do;
> ecd[n]:=eval(Y,t=c0)=alpha:
> ecd[n+1]:=eval(Y,t=d0)=beta:
> C:=solve({seq(ecd[k],k=0..n+1)},
> [seq(c[k],k=0..n+1)]);
> assign(C):
> return Y:
> end proc:

```

The mesh points are the roots of the  $n$ -th degree second kind Chebyshev polynomials (formula (7)) and the points  $c$  and  $d$ .

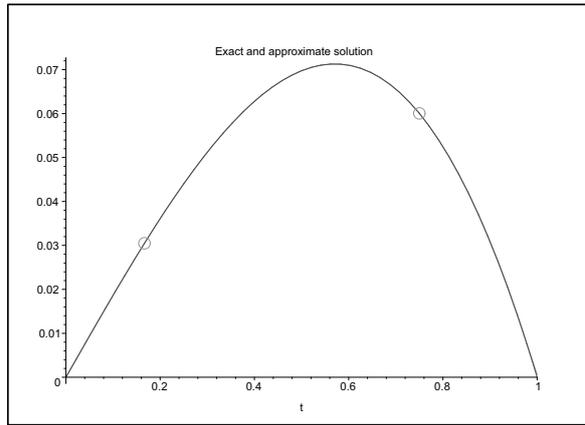
#### 4 Numerical examples

We present two examples: one with a nonoscillating solution and the other with oscillating solution. A problem with a nonoscillating solution is simple and does not require a large computational effort. A problem with an oscillating solution is harder, and requires a mesh with a large number of points. The methods do not depend on conditions on  $q(x)$ . We solved our examples using both methods. For each example and method we plot the exact and the approximate solution and generate the execution profile (with the pair `profile - showprofile`). The first example is from [3, page 560]

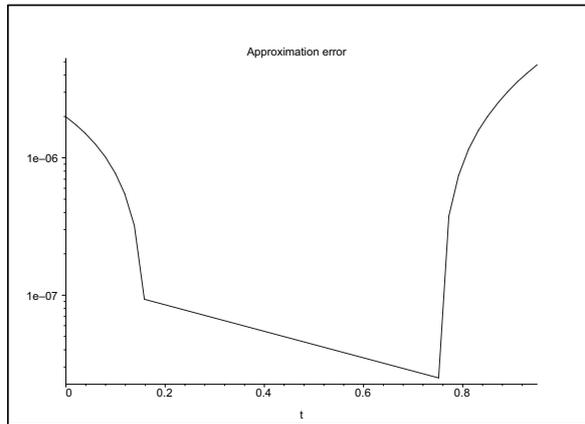
$$\begin{aligned}
 -y'' - y &= x, & x \in [0, 1] \\
 y\left(\frac{1}{6}\right) &= -\frac{1}{6} \frac{-6 \sin \frac{1}{6} + \sin 1}{\sin 1}, \\
 y\left(\frac{3}{4}\right) &= -\frac{1}{4} \frac{-4 \sin \frac{3}{4} + 3 \sin 1}{\sin 1}.
 \end{aligned}$$

The exact solution is  $Z(t) = -\frac{-\sin(t)+t \sin 1}{\sin 1}$ , and we computed it using `dsolve`. We chose  $n = 10$  for both methods and  $k = 3$  for the first method. Figure 1 shows the exact and the approximate solution computed using the first method. The error plot in a semilogarithmic scale is given in Figure 2.

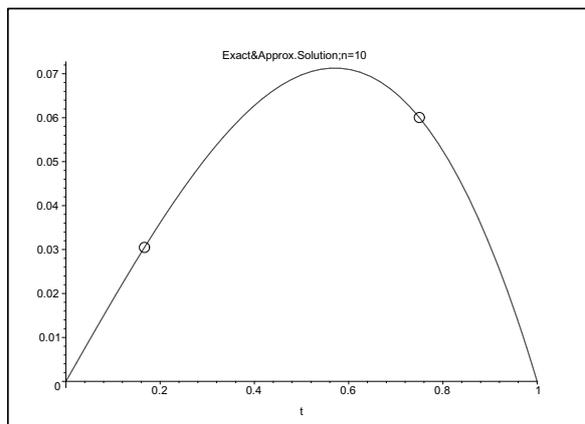
The corresponding graphs for Chebyshev methods are illustrated in Figures 3 and 4.



**Figure 1.** The graph of exact and approximate solution, nonoscillating problem, B-spline method,  $n = 10, k = 3$



**Figure 2.** Error plot, nonoscillating problem, B-spline method,  $n = 10, k = 3$



**Figure 3.** Exact and approximate solution, nonoscillating problem, Chebyshev method,  $n = 10$

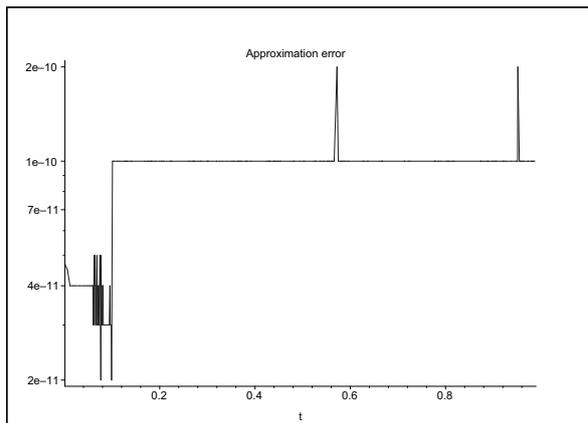


Figure 4. Error plot, nonoscillating problem, Chebyshev method,  $n = 10$

Here are the profiles for the procedures in the case of nonoscillating problem. The function showprofile for the B-spline method gives the following results:

function	depth	calls	time	time%	bytes	bytes%
genspline	1	1	2.496	100.00	92768440	100.00
total:	1	1	2.496	100.00	92768440	100.00

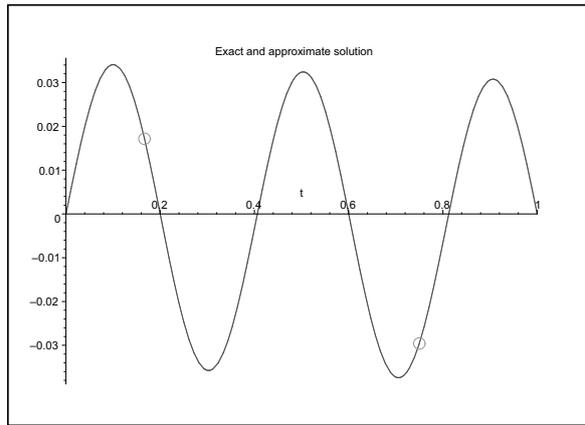
The profile for Chebyshev method is:

function	depth	calls	time	time%	bytes	bytes%
genceb	1	1	0.249	100.00	9291004	100.00
total:	1	1	0.249	100.00	9291004	100.00

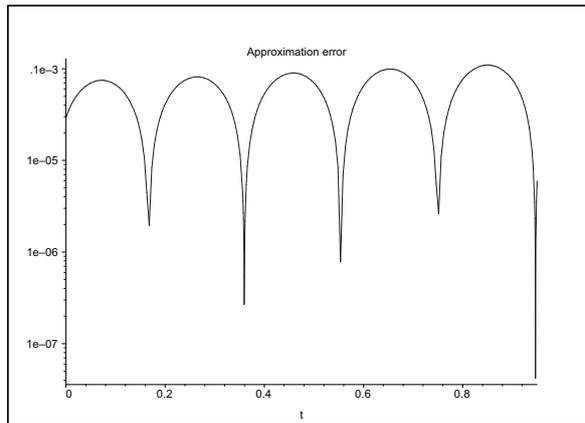
The second example has an oscillating solution:

$$\begin{aligned}
 -y'' - 243y &= x, \quad x \in [0, 1] \\
 y\left(\frac{1}{6}\right) &= -\frac{1}{1458} \frac{-6 \sin \frac{3}{2}\sqrt{3} + \sin 9\sqrt{3}}{\sin 9\sqrt{3}} \\
 y\left(\frac{3}{4}\right) &= -\frac{1}{972} \frac{-4 \sin \frac{27}{4}\sqrt{3} + 3 \sin 9\sqrt{3}}{\sin 9\sqrt{3}}.
 \end{aligned}$$

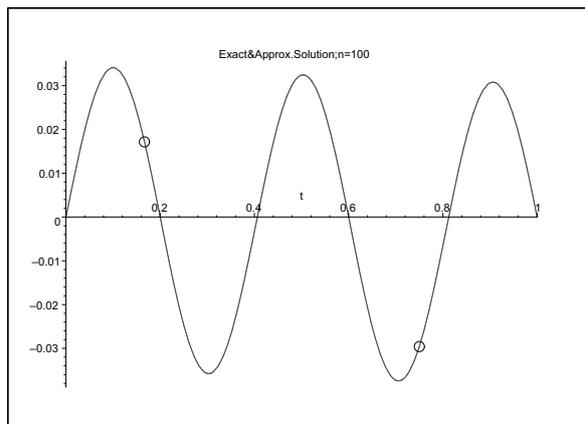
The exact solution, provided by dsolve is  $Z(t) = -\frac{1}{243} \frac{-\sin 9\sqrt{3}t + t \sin 9\sqrt{3}}{\sin 9\sqrt{3}}$ . We chose  $n = 100$  for both methods and  $k = 3$  for the first method. Figure 5 gives the graph of exact and approximate solution for the oscillating problem. The error plot appear in Figure 6. The corresponding graphs for Chebyshev methods are given in Figures 7 and 8, respectively.



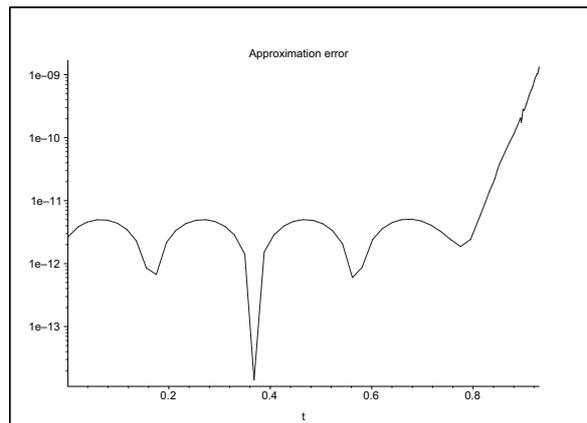
**Figure 5.** Exact and approximate solution, oscillating problem, B-spline method,  $n = 100$ ,  $k = 3$



**Figure 6.** Error plot, oscillating problem, B-spline method,  $n = 100$ ,  $k = 3$



**Figure 7.** Exact and approximate solution, oscillating problem, Chebyshev method,  $n = 100$



**Figure 8. Error plot, oscillating problem, Chebyshev method,  $n = 100$**

Here are the profiles for the procedures in the case of oscillating problem.

function	depth	calls	time	time%	bytes	bytes%
genspline	1	1	36.691	100.00	990061444	100.00
total:	1	1	36.691	100.00	990061444	100.00
function	depth	calls	time	time%	bytes	bytes%
genceb	1	1	174.814	100.00	4365866236	100.00
total:	1	1	174.814	100.00	4365866236	100.00

## 5 Conclusions

The Chebyshev method has a smaller error (see error plots, Figures 2, 4, 6, 8). For the nonoscillating solution and a mesh with a small number of subintervals Chebyshev method is faster and requires less memory. If the number of points increases the B-spline method is faster and requires less memory. The reason is that for the B-spline method the matrix of the system that provides the coefficients is a band matrix with at most 4 nonzero elements per line, while for Chebyshev method the matrix is dense. The example with oscillating solution supports this conclusion.

Our approach based on computer algebra has the following advantages:

- The choice of mesh points is arbitrary.
- The degree of Legendre polynomial can be changed.
- We need not bother with differentiation, equation building, ordering and so on.
- The analytic form of the solution allow to compute the approximation at any point, to plot it and to use it further as input for other problems.

## References

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